

HW 2

1. (a) sign

$$S = M^d \pmod{33} \\ = 25^7 \pmod{33} = \cancel{19} 31$$

(b) verify

$$M' = S^e \pmod{33} \\ = 31^3 \pmod{33} \\ = 25 = M$$

2. $p=23, d=9, g=5, X=10$

Select $K_E = 7 \in \mathbb{Z}_{22}$ s.t. $\gcd(7, 22)$

$$K_E^{-1} = 19 \pmod{22}$$

sign

$$r = 5^7 \pmod{23} \\ = \boxed{17} \pmod{23}$$

$$S = (10 - 9 \cdot 17) \cdot 7^{-1} \pmod{22} \\ = \boxed{11} \pmod{22}$$

Alice $\xrightarrow{(10, (17, 22))}$ Bob

verify

$$V = \beta^r r^S \pmod{p} \\ = \cancel{17}^9 \cdot 17^{11} \pmod{23} \\ = 5^{10} \pmod{23}$$

$$\left(\begin{aligned} \beta &= g^d \pmod{p} \\ &= 5^9 \pmod{23} \\ &= 11 \pmod{23} \end{aligned} \right)$$

$$V' = g^X \pmod{p}$$

$$= 5^{10} \pmod{23}$$

$$\therefore V = V'$$

3. $p=27, q=13, X=5, H(x)=5$

choose $g=6$ s.t. $g^q = 1 \pmod{p}$
($\gcd(g, p) = 1$)

choose $d=7$

(a) select $K=2, K^{-1} \pmod{q} = 2^{-1} \pmod{13} = 7$

$$\begin{aligned} r &= (g^K \pmod{p}) \pmod{q} \\ &= (6^2 \pmod{27}) \pmod{13} \\ &= \boxed{9} \pmod{13} \end{aligned}$$

$\frac{36}{27} = 9$

$$\begin{aligned} s &= K^{-1} (H + d \cdot r) \pmod{q} \\ &= 2^{-1} (5 + 7 \cdot 9) \pmod{13} \\ &= 7 \cdot 68 \pmod{13} \\ &= 7 \cdot 3 \pmod{13} \\ &= \boxed{8} \end{aligned}$$

(b) Verify

$$\begin{aligned} s' &= 8^{-1} \pmod{13} = \boxed{5} \pmod{13} \\ (8 \cdot 8^{-1} \pmod{13} &= 1 \pmod{13}) \end{aligned}$$

$$\begin{aligned} r' &= (\beta^{s' r} g^{s' H'} \pmod{p}) \pmod{q} \quad (\beta = 6 \pmod{p}) \\ &= (\beta^{5 \times 9} 6^{5 \times 5} \pmod{27}) \pmod{13} \end{aligned}$$

check $r' \stackrel{?}{=} 9$

4. (a) In RSA, $S = M^d \pmod n$
 (M: message)

So, S can be as large as $(n-1)$
 the size of s, $|s| \approx |n| \approx 1024$ bits

(b) In ElGamal, S ,

$$r = \square \pmod p, \quad s = \square \pmod{p-1}$$

So, $|r| = p-1 \approx 1024$ bits, $|s| = |p-2| \approx 1024$ bits

the size of signature is 2048 bits.

(c) In the DSS,

$$r = \square \pmod q, \quad s = \square \pmod q$$

$$|r| = |s| = |q-1| \approx 160 \text{ bits}$$

So, the size of signature is 320 bits.

6. input: $(Tlen, \underset{\substack{\uparrow \\ \text{AES key}}}{K}, k_i, \underbrace{M = (M_1, \dots, M_N)}_{M[1], \dots, N}, N \text{ blocks})$

$$\{ C[i] \leftarrow E_K(M[i])$$

$$i \leftarrow 2$$

while $(i < N)$

$$C[i] \leftarrow E_K(C[i-1] \oplus M[i])$$

$$i \leftarrow i + 1$$

$$C[N] \leftarrow E_K(C[N-1] \oplus M[N] \oplus k)$$

$$h \leftarrow \text{selectLeft}(C[N], Tlen)$$

return h.

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